Logic Sniffer Specification
(Demon Core FPGA)

Version 1.01

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Table of Contents

Table of Figures

1   Introduction

The Open Logic Sniffer Demon Core FPGA offers many of the features found in an HP 16500 / 16550 timing logic analyzer.   Combined with the OLS's small size it yields a very potent little tool.

Feature list:

- SUMP logic analyzer client compatible.
- 16 pin buffered probe header (5V tolerant)
- 16 pin unbuffered (wing) header (3.3V tolerant)
- 24K samples (8 bit captures)
- Maximum 200Mhz sample rate (using demux mode)
- USB interface, USB powered, & USB upgradeable.
- Basic/Legacy trigger with four level parallel or serial target matching
- Advanced trigger with 16 sequence states, 10 pattern matches, two range checkers, two edge detectors & two timers (10ns resolution, 600 second range).
- RLE compression of captures.

The Logic Sniffer was developed as a collaboration between the Gadget Factory (http://www.gadgetfactory.net) & Dangerous Prototypes.(http://dangerousprototypes.com).

Figure 1 - Block Diagram

1.1 Advanced Trigger

The advanced trigger offers many features found in an HP16500 / 16550 timing logic analyzer. It offers 16 sequence levels with two way branching (a programmable state machine), ten pattern value comparisons, two range checks, two edge checks, and two counter/timers.

There are several different modules within the advanced trigger, each needing its own spin on configuration. These are:

- Trigger Terms. Terms perform the actual bitwise masked pattern comparisons on the 32-bit incoming data. There are ten terms. Each one needs 128 bits of configuration data.

- Trigger Range Detectors. The two range detectors look for values falling between a lower & upper 32-bit limit. Each range check uses two 32-bit magnitude comparisons., one for the upper compare, one for the lower. A match means: "upper >= indata >= lower". Each range detector needs 512 bits of configuration data.

- Trigger Edge Detectors. The two edge detectors compare sampled input against a delayed version. Rising edges, falling edges, or both, or neither can thus be detected. Each needs 256 bits of configuration data.

- Trigger Timers. Timers are started/stopped under control of the trigger sequence states. There are two timers, each needing a 36 bit limit value. The timers have a range from 10ns to over 11 minutes.

- Trigger Sums. These combine the results of the trigger terms, range comparisons, edge checks, and timer checks. There are 11 logic units programmable to perform any logic function. Examples are AND/NAND/OR/NOR/XOR/NXOR. The logic units are combined in a tree structure, and combine/sum all trigger sources down to a match/miss signal. There are three trigger sums per trigger state, and 16 trigger states, totaling 48 sums. Each sum needs 192 bits of configuration data.

- Trigger Sequence States. Each state is guided by the results of three trigger sums (see above). If the "hit" trigger sum matches, for a specified number of occurrences, the FSM advances. A hit also controls if a timer is started/stopped/cleared. If the "else" trigger sum matches, the FSM branches to the "else" state. Data capture is controlled by the "capture" trigger sum – if it matches, a sample is taken. If not, then not. Each state takes 32 bits of configuration.

Each module will be described in detail in the following chapters.

2 Background

A little background to help put things in context...

2.1 FPGA's

The FPGA (field programmable gate array) is very useful for creating custom hardware. They can implement anything from simple logic up to entire computers. To accomplish this, an fpga contains large numbers of tiny lookup tables and flip-flops as building blocks, plus usually some larger memories (SRAM's).
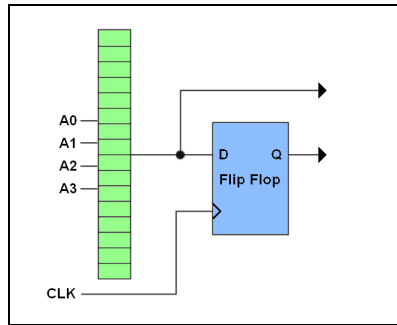
Figure 2 - Logic Block

Add some I/O pins for talking to the outside world, plus lots of configurable wires to connect everything, and you have an FPGA.

The lookup tables are the interesting piece, and store only 16 bits of data. They have four inputs (the address), an output, and some config pins. Normally they are configured to form logic. Any four-input function can squeeze in there.

For example, a four input AND gate. Load LUT addresses 0 to 14 with all zeros, and set address 15 to one. Now if all inputs are set, you get a one on output -- otherwise zero. Just like an AND gate.

Figure 3 - FPGA Layout
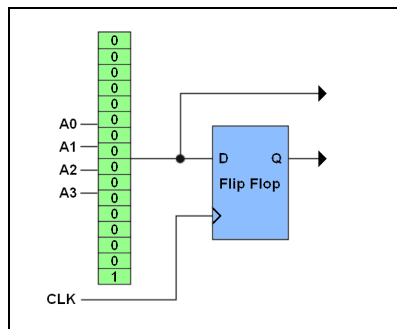
Figure 4 - AND gate equivalent

In the case of a logic analyzer, even more useful functions are possible. For example, a trigger search on four bits of sampled data. The LUT allows easy matching of any pattern.

LUT's are normally programmed serially when the FPGA powers up. Think shift register. However, they can also be selectively changed during runtime.

2.2   The Datapath

The Logic Sniffer is composed of several main blocks.   Most are concerned with the data path (see right).   In the drawing thin lines are control signals.   Thick lines are the data path.

The "*sync module*" captures input data using the sample clock. The sample clock can either be the internal 100Mhz reference clock, or supplied from an external source.

The "*async fifo*" transfers the captured input to the rest of the fpga. If the sample clock & internal reference clock are totally unrelated, the fifo avoids timing problems.

The "*sampler module*" is what implements the "sampling rate". Assume the sample clock is using the 100Mhz reference.   If you want a 10Mhz sample rate, the sampler forwards every tenth valid capture.   For 1Mhz, every hundredth valid capture, etc…

The "*triggers*" look for patterns. They control what ultimately gets stored in SRAM -- with the "*capture*" signals -- and tell the controller when a trigger hit occurs.   The basic trigger captures everything always.   The advanced trigger lets you be more selective.

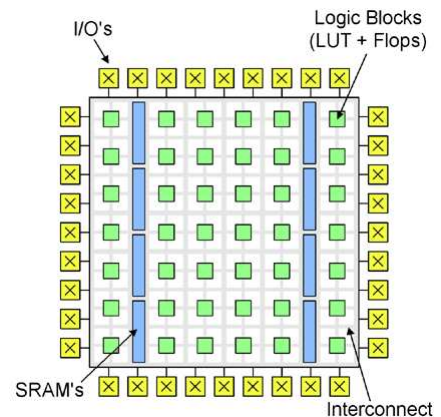The "*delay fifo*" aligns its output to the capture/run trigger outputs. The trigger modules take a few clocks to evaluate things.   For a trigger "capture" to work, the input to the next state must match.

The "*data align*" module removes gaps within the input data caused by disabled groups. If you decide to disable groups 0 & 2, the aligner shifts the remaining groups down to fill the gap.   This simplifies things later.

The "*RLE encoder*" looks for repeat captures.   If found, counts are stored instead of repeats of unchanging data.   This can greatly increase the effective storage capacity.

The "*controller*" is what decides how much gets captured, and when to send results to the client on your PC.

The "*SRAM interface*" handles saving & fetching captured data to & from the SRAM memory.

The "*SPI interface*" uploads the captures serially to the client. Only valid data is sent.   Gaps due to disabled groups are filtered out.

A few remaining blocks aren't shown here.   The inbound SPI interface for receiving commands, the command decoder & flags register.
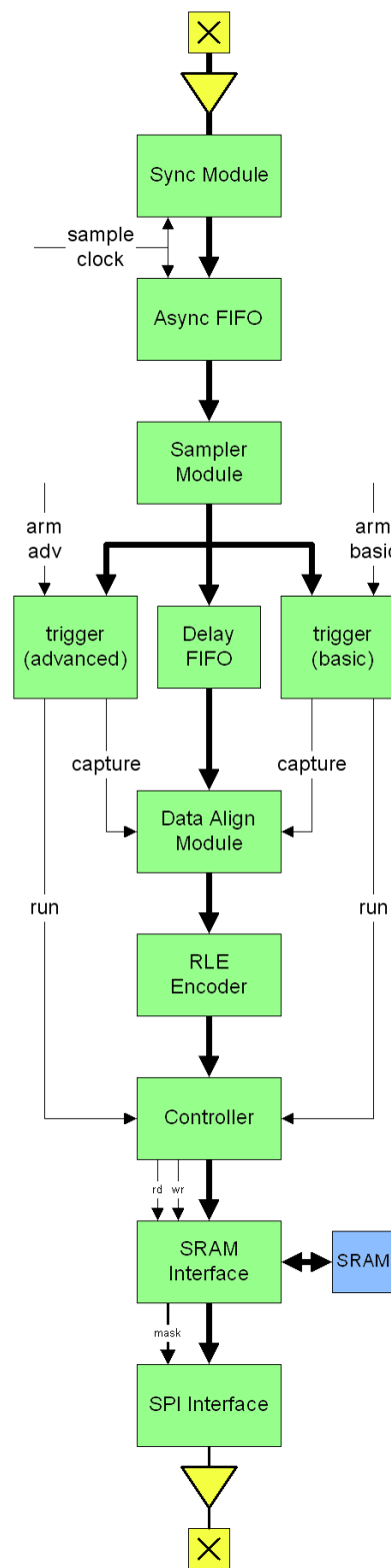
Figure 5 - Datapath Diagram

2.3 The SUMP Protocol

The SUMP communication protocol assumes a standard RS232 port connection.   Although the Logic Sniffer is USB based it emulates a very fast serial port from the perspective of the PC.

When sending captured data the analyzer sends back a contiguous block of data.   Disabled channels are not sent however.   First byte sent corresponds to sampled bits 7:0 (assuming channel 0 isn't disabled).

The following list documents the commands understood by the Logic Sniffer FPGA.

2.3.1 Short Commands (one byte commands)

0x00 - Reset

Reset the analyzer.   Send at least five times to ensure FPGA in a known state.   Clears any previous long commands that might have gotten stuck.

0x01 - Arm Basic/Legacy Trigger

Arm the basic trigger.   Sampling begins immediately.   Once the trigger fires, the controller waits for "delay count" additional samples before returning captured data to client.

0x02 - Query ID

The device responds with four bytes.   Currently "SLA1" to maintain backwards compatible SUMP protocol, output LSB first. ie: "1", "A", "L", "S"

0x04 - Query Meta Data

The device responds with a variable length block of data composed of one or more tokens + associated data.

- Token 0 denotes end of meta data
- Tokens 0x01 to 0x1F are followed by a null terminated string (UTF-8 encoded).
- Tokens 0x20 to 0x3F are followed by a 32-bit unsigned integer (MSB first).
- Tokens 0x40 to 0x5F are followed by a 8-bit unsigned integer.
- All other token ranges are reserved.

The following fields are currently defined:

Token	Description	FPGA Currently Returns
0x00	End of Meta Data	0
0x01	Device Name	"Open Logic Sniffer v1.01"
0x02	Version of FPGA firmware	"3.0"
0x03	Ancillary Version (PIC firmware)	-
0x20	Number of Probes	-
0x21	Sample memory available (bytes)	24576
0x22	Dynamic memory available (bytes)	-
0x23	Maximum sample rate (Hz)	200000000
0x24	Protocol version	-
0x40	Number of Probes (short version)	32
0x41	Protocol version (short version)	2

Clients should accept data from either long or short versions of a meta field, if both are possible.

0x05 - Finish Now (disable RLE mode)

Turns off RLE mode specified in the flags register (see 0x82 command below).   In RLE mode, the analyzer can sit forever if inputs aren't changing.   This gives the PC client application an opportunity to finish a capture on demand.

Finish now also aborts the advanced trigger, & forces data capture to quickly fill the buffer & end capture.

0x06 - Query Input Data

The device responds with four bytes. Gives a snapshot of the current logic analyzer input bits.

0x0F - Arm Advanced Trigger

Arm the advanced trigger.   Conditional capturing (under control of trigger sequencer) begins immediately.   Once trigger fires, the controller waits for "delay count" additional samples before returning captured data to client.

0x11 - XON

Depreciated.   On a serial port, this would pause output when FPGA was returning captured data.   Currently ignored by the FPGA.

0x13 - XOFF

Depreciated.   On a serial port, this would resume output when FPGA was returning captured data. Currently ignored by the FPGA.

2.3.2 *Long Commands (five byte commands)*

0x80 - Set Sample Rate

The Logic Sniffer samples at rates lower than the reference clock by using a *divisor-count*.   Every given number of clocks, a sample is accepted.   To set the sample rate, the divisor count is computed as follows:

$$\text{divider} = (\text{clock} / \text{sample-rate}) - 1$$

0x81 - Set Read Count & Delay Count

The *read-count* specifies the number of samples (divided by four) to be returned to the client on your PC.   The *delay-count* is the number of samples (divided by four) to capture -after- the basic or advanced trigger fires.

A *read-count* larger than the *delay-count* returns captures from before the trigger.   Such pre-trigger captures are only valid if the device was running long enough before the trigger fired.   ie: a trigger that fires immediately could return garbage.

0x82 - Set Flags

The flags register configures a whole range of miscellaneous Logic Sniffer features including:

- Demux Mode. Double-Data-Rate Capturing. Not compatible with filter mode. Should be used with max sample-rate (ie: set divisor to zero -- see 0x80 above). 1= Enabled.
- Noise Filter Mode. Reduces occurrence of glitches on inputs. 1 = Enabled.
- Disabled Channel Groups. Don't capture specified groups of input data. 1 = Disable group.
- External Clock Source. Use external clock for sampling. If disabled, use internal 100Mhz reference clock instead. When enabled the Logic Sniffer is basically in "state mode" (sample synchronous with input) verses "timing mode" (asynchronous internal reference clock). 1 = Enabled.
- Inverted External Capture Clock. Capture data on falling edge of sample clock when using external clock.
- Enable RLE compression mode. Duplicate samples are compressed into counts. Most significant bit of sample replaced with "rle-flag", and is thus no longer available. Requires client support to decode. A value is stored, followed by a count of repeats. Count's are exclusive of value, so <value><count=7> means there are 8 samples.
- Swap Number Scheme. Swap upper 16 bits & lower 16 bits of input in capture buffer. 1 = Enabled.
- External test mode. Output 16-bit test pattern on capture pins 31:16. 1 = Enabled.
- Internal test mode. Supply internally generated 32-bit test pattern to capture buffer. Input data is ignored.
- Extra RLE compression modes. For input samples which change infrequently, rle-count's can fill to maximum. The Logic Sniffer then stores that count & starts counting again. The extra modes control how often another <value> is stored (even if it hasn't changed):

| Mode	Description
0&1	Issue <value> & <rle-count> as pairs. Backwards compatible.
2	Periodic. <values> reissued approximately every 256 <rle-count> fields.
3	Unlimited. <values> can be followed by unlimited numbers of <rle-counts>.

0x9E - Advanced Trigger Configuration Select

The advanced trigger needs a lot of configuring (84 different selectors), so it seems prudent to provide a separate address space to cover programming the 11000+ config bits.   This register selects what to configure in the advanced trigger.   Please see Appendix A for the list of advanced trigger registers.

0x9F - Advanced Trigger Write Data Register

The *trigger-data* register loads the LUT chains of trigger terms, range checks, edge detects & summing terms.   It is also used to write the sequence states & timer limit registers.

0xC0/C4/C8/CC - Set Basic Trigger Mask

The *trigger-mask* controls which basic *trigger-value* bits participate in a compare.   A mask bit set to zero causes the corresponding trigger-value bit to be ignored.   If the entire mask is zero, the trigger will match everything always.

0xC1/C5/C9/CD - Set Basic Trigger Value

The *target-value* is what the fpga looks for in the sampled input data.   Each bit corresponds directly 1:1 with the captured input data.   Note: Disabled groups still participate in a trigger.   If you want to exclude a disabled group, then be sure to set the trigger mask accordingly.

NOTE: The FPGA requires trigger masks & values be written back-to-back. ie: Write the mask, then the value, for trigger 0. Then the mask/value for trigger 1, etc...   The basic triggers use the same LUT based compare mechanism as the advanced triggers.   For backwards compatibility, a widget converts the mask/value pairs into serial LUT chains automatically for the basic trigger.

0xC2/C6/CA/CE - Set Basic Trigger Configuration

Writing this register configures the selected basic trigger stage. There are various fields available for customizing the operation of the trigger:

- Stage Delay. If a trigger match occurs, the action of the stage is delayed by number of samples.
- Serial Mode. If enabled, trigger operates as a serial trigger. 1 = Serial Mode.
- Serial Channel. In serial mode, select which input pin to treat as serial channel (0 to 31).
- Level. Trigger level at which current stage will become active. Level 0 means activate immediately.
- Start. Tell controller the basic trigger has fired.
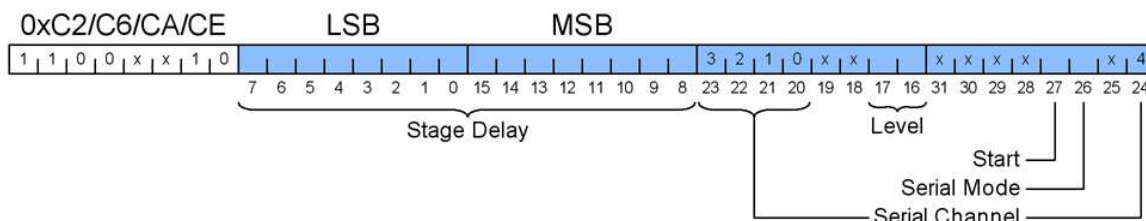
Serial Mode:

In serial mode, a trigger stage monitors only a single sampled input bit, selected with the "*serial-channel*". The selected bit is fed into a serial-to-parallel shift register. Thus the last 32 samples of a single bit can be evaluated using the normal masked compare logic.

To be effective, you should use an external clock in serial mode. Otherwise it's very difficult to predict exactly when samples are taken & added to the shift register. So for example, when monitoring a SPI bus use the SPI clock for capturing.
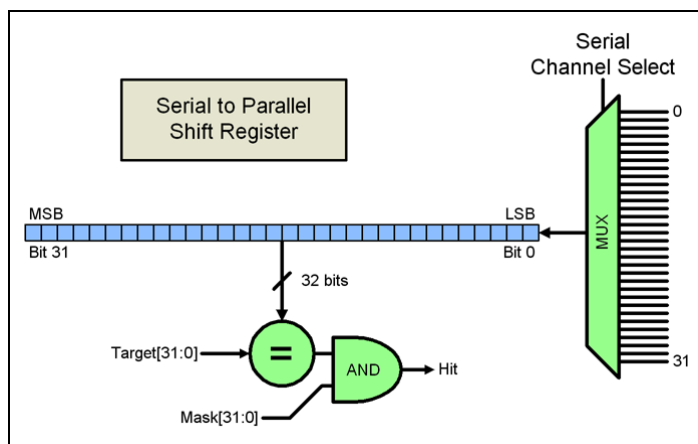

Figure 6 - Basic Trigger Serial Mode

Trigger Levels:

The basic trigger has a simple level counter. When any trigger stage matches something, the level count is incremented. When the level count matches a given stages "*trigger-level*", the stage becomes active. Thus you can search for a simple sequence of events.

ie: Setup trigger 0 to start immediately, triggers 1 & 2 to start on level 1, and trigger 3 to start on level 2. When trigger 0 matches, both triggers 1 & 2 activate. When either of them matches, trigger 3 activates.

3   *Trigger Terms*

A trigger term is pretty simple.   It compares the logic analyzer input against a target.   If the input matches the target, it signals a hit.   Trigger terms work the same in both basic & advanced triggers.

To make things a little more useful, the fpga also supports masking (or ignoring) parts of the input.   The mask determines which bits must match, so you only compare them & nothing else.

As shown below, the compare is done with an XNOR (exclusive inverted or), which goes high when both inputs match.   Masking is done by an AND gate.
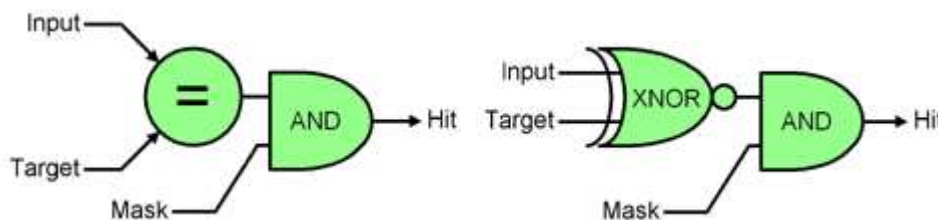
Figure 7 - Masked Compare

To implement something like this normally requires flops to hold the target, flops to hold the mask, and logic for the XNOR & AND gates. So around 64 flops plus 32 logic blocks.

However, a better way exists.   In an fpga, LUT's can handle everything. The 32-bit target & mask values, and logic are all encoded directly into only eight LUT's.

Quite a savings, but there is a downside. Programming the LUT is non-trivial.

The "complexity" of the masked compare doesn't go away -- its simply transferred elsewhere.   Thus in the advanced trigger, the eight LUT's need 128 bits of configuration supplied by the client.

The LUT memory is programmed with the -results- of a masked compare.   The target & mask are known, and the input sample is one of sixteen possible values.

Thus we simply run through all combinations in advance. ie:

Mask	Input	Target	Hit
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

Figure 8 - Truth Table

Figure 9 - Trigger Term LUT

```
LUT address 0 stores result of:  (0 ^ value) & mask)==0
LUT address 1 stores result of:  (1 ^ value) & mask)==0
LUT address 2 stores result of:  (2 ^ value) & mask)==0
LUT address 3 stores result of:  (3 ^ value) & mask)==0
LUT address 4 stores result of:  (4 ^ value) & mask)==0
etc…
```

For each of the eight LUT's, the above is repeated with different 4 bit chunks of value & mask.
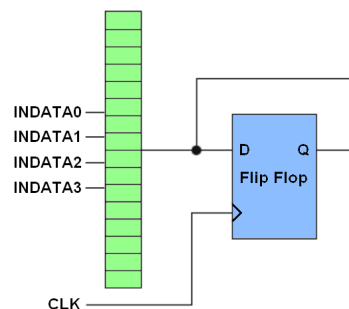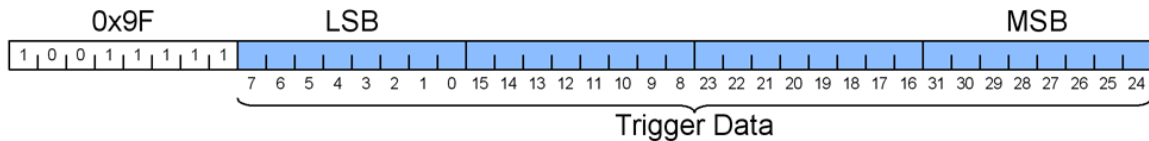
3.1 Programming Lookup Tables

To enable easier LUT programming, the fpga contains a built in parallel-to-serial converter.   You simply write a long SUMP command to the chain address.   Note the SUMP protocol dictates the least significant data byte be output first.

Once the fpga receives the long command, the bytes are organized MSB first (bits 31 downto 0), and then serially shifted into the LUT chain.
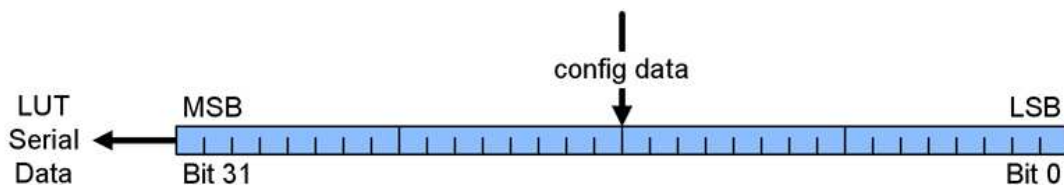
Some of the advanced trigger modules are quite complex and need multiple writes to fully load their LUT chains. Chains vary in length up to 512 bits.

3.2 Example (Trigger Term Initialization)

```
#define write_select(value) write_long_command (0x9E, value)
#define write_chain(value) write_long_command (0x9F, value)

void write_term (int term_number, unsigned int target, unsigned int mask)
{
  unsigned int bitmask=1;
  unsigned int lutvalue0=0;
  unsigned int lutvalue1=0;
  unsigned int lutvalue2=0;
  unsigned int lutvalue3=0;

  for (int i=0; i<16; i=i+1) {
    if (((i ^ (target & 0xF)) & (mask & 0xF))==0) lutvalue0 |= bitmask;
    if (((i ^ ((target>>4) & 0xF)) & ((mask>>4) & 0xF))==0) lutvalue0 |= (bitmask<<16);
    if (((i ^ ((target>>8) & 0xF)) & ((mask>>8) & 0xF))==0) lutvalue1 |= bitmask;
    if (((i ^ ((target>>12) & 0xF)) & ((mask>>12) & 0xF))==0) lutvalue1 |= (bitmask<<16);
    if (((i ^ ((target>>16) & 0xF)) & ((mask>>16) & 0xF))==0) lutvalue2 |= bitmask;
    if (((i ^ ((target>>20) & 0xF)) & ((mask>>20) & 0xF))==0) lutvalue2 |= (bitmask<<16);
    if (((i ^ ((target>>24) & 0xF)) & ((mask>>24) & 0xF))==0) lutvalue3 |= bitmask;
    if (((i ^ ((target>>28) & 0xF)) & ((mask>>28) & 0xF))==0) lutvalue3 |= (bitmask<<16);
    bitmask <<= 1;
  }

  // Write data into LUT serial chain.  MSB must goes in first.  Total of 128 bits.
  write_select (0x20 + (term_number%10));
  write_chain (lutvalue3);
  write_chain (lutvalue2);
  write_chain (lutvalue1);
  write_chain (lutvalue0);
}
```

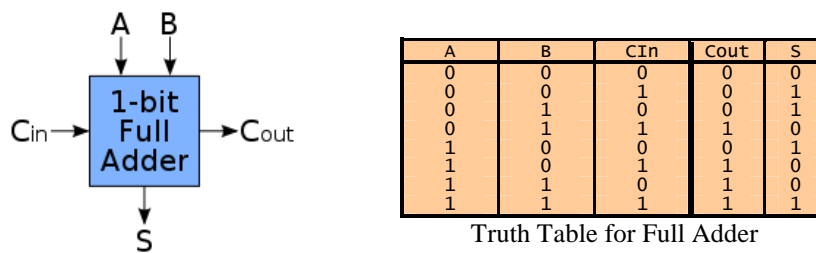**Figure 10 - Initialize Trigger Term LUT Values**

4   Range Detectors

The advanced trigger supports detecting when sampled data falls within a range of values.   ie: if the sample is between a lower limit & upper limit. ie:

$$hit = (upper\text{-}limit >= sample)\ \&\&\ (sample >= lower\text{-}limit)$$

So we need magnitude comparators.   In hardware, that means adding or subtracting two numbers to see if something overflows or "carry's" over.   A bit like long addition.   Add a column, carry the one, add another, etc...

In hardware there is parallel to long addition: the full adder.   The adder adds A+B+Cin, and outputs S & Cout, where all signals are single bits.   To add wide values you chain the carry signals as needed.

A	B	CIn	Cout	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

Truth Table for Full Adder

Chains of full adders work, but are slow.   The carry-out signal is derived from the inputs.   In a 32-bit adder, there would be 32 chained carry circuits.   Achieving timing would be harder than necessary.

An improvement is the carry-look-ahead adder.   There is still a combinatorial element per bit, but it's very fast -- a simple mux.   The fpga even has an integral "fast-carry-chain" circuit available for implementing such adders.
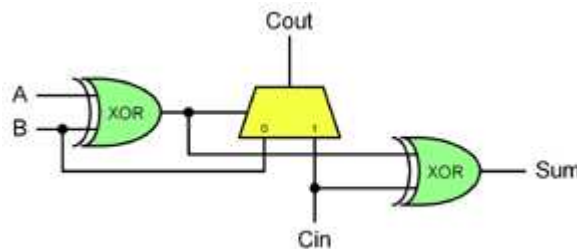

Figure 11 - Carry Look Ahead Adder

To make things even better, a magnitude comparison doesn't need the sum output at all.   Just the fast carry chain.   Thus we end up with something that fits perfectly into a single Logic Block.   The LUT table simply needs loading with the result of an XOR between the input on A0, and the target value.
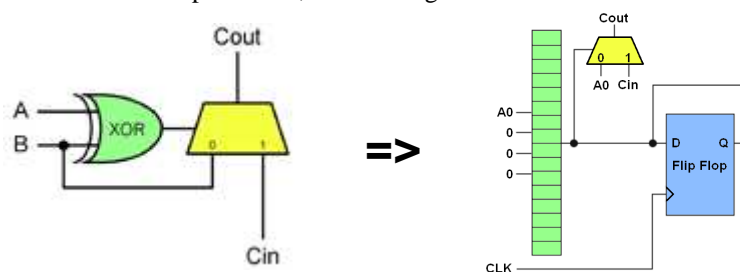

Figure 12 - Carry Look Ahead Comparator

4.1 *Complex Range Compares*

The fast carry change logic has a couple of interesting possibilities.

First of all, the range comparison can start anywhere.   If the first bit of interest is bit 7 (for example), simply fill all lower LUT's with 1's.   This NOP's them in the comparator.

Second, the bits participating in the range comparison don't need to be contiguous.   You could setup to range check the value formed by indata bits 7,9,11, & 15.   Simply NOP the LUT's of the don't care indata bits.

4.2 *Prepare Target Values for Comparator*

There are two ways to compare a value.   Subtract a value & look for underflow/borrow, or add & look for overflow/carry.   Since we have the fast-carry-chain, we're using the latter approach.

A given range target must be bitwise inverted before being programmed.   For non-contiguous range checks, the inverted value must also be spaced out, across only those bits of interest.

> Lower value = ~(lower_target - 1)
> Upper value = ~(upper_target)

Yields a hit if (upper_target >= indata >= lower_target)

4.3 *Range Compare LUT Chains*

There are 32 LUT's in each range-check magnitude comparator.   Each needs 16-bits, so that totals 512 bits of configuration.   Two LUT's are programmed at once for each write to the chain-data register.
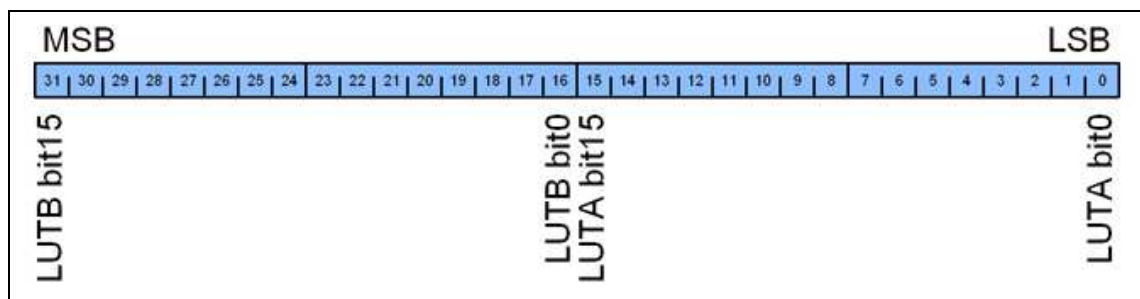
Figure 13 - Range LUT Organization

4.4 *Example (Range Detect Initialization)*

```
#define RANGE_XOR0 0xAAAA
#define RANGE_XOR1 0x5555
#define RANGE_NOP 0xFFFF

//
// Setup LUT's for Range Detectors
//
// Inputs:
//   rangesel:  0=range1-lower, 1=range1-upper, 2=range2-lower, 3=range2-upper
//   target:    Limit value.
//   mask:      Indicate which bits should participate in range compare.
//
void write_range (int rangesel, unsigned int target, unsigned int mask)
{
  unsigned int value;
  unsigned int lutvalue=0;
  int i;

  write_select(0x30+(rangesel&3));

  // Count # of bits in mask...
  int bitcount=0;
  unsigned int bitmask=1;
  for (i=0; i<32; i++, bitmask<<=1)
    if (mask & bitmask)
      bitcount++;

  // Prepare target value...
  if (rangesel & 1)
    value = ~target; // upper range target
  else value = ~(target-1); // lower range value

  // Push MSB of target into bit 31...
  value <<= (32-bitcount);

  // Generate & program LUT values.  Total of 512 bits.
  for (i=0; i<16; i=i+1) {
    if (((mask>>31)&1)==0)
      lutvalue = RANGE_NOP;
    else {
      lutvalue = ((value>>31)&1) ? RANGE_XOR1 : RANGE_XOR0;
      value <<= 1;
    }
    mask <<= 1;
    lutvalue <<= 16;

    if (((mask>>31)&1)==0)
      lutvalue |= RANGE_NOP;
    else {
      lutvalue |= ((value>>31)&1) ? RANGE_XOR1 : RANGE_XOR0;
      value <<= 1;
    }
    mask <<= 1;

    write_chain(lutvalue);
  }
}
```

Figure 14 - Initialize Range Detector LUT Values

5   Edge Detectors

The edge detectors compare a delayed sample to the current one.   If any bit changes, with a rising edge, falling edge, either rising or falling edge, or neither (no change), it can be detected.

Each 4-input logic block evaluates two bits of input, and two bits of delayed input.
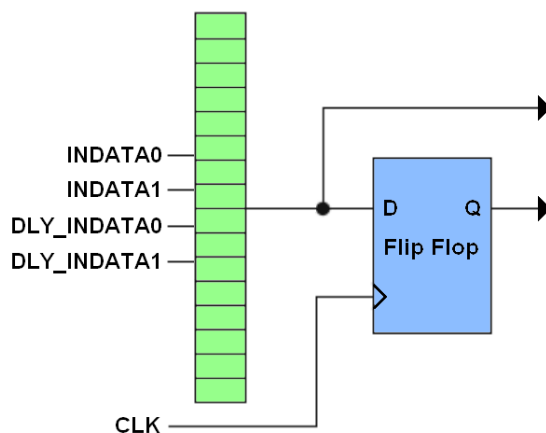
Figure 15 - Edge Detector LUT

All that is left now is placing 1's in the LUT where the edges of interest occur.   Since two bits are evaluated in parallel, this becomes slightly more complex.

The following truth table shows the possible combinations.

DLY INDATA1	DLY INDATA0	INDATA1	INDATA0	RISE0	FALL0	RISE1	FALL1	BOTH0	BOTH1	NONE0	NONE1
0	0	0	0	0	0	0	0	0	0	1	1
0	0	0	1	1	0	0	0	1	0	0	1
0	0	1	0	0	0	1	0	0	1	1	0
0	0	1	1	1	0	1	0	1	1	0	0
0	1	0	0	0	1	0	0	1	0	0	1
0	1	0	1	0	0	0	0	0	0	1	1
0	1	1	0	0	1	1	0	1	1	0	0
0	1	1	1	0	0	1	0	0	1	1	0
1	0	0	0	0	0	0	1	0	1	1	0
1	0	0	1	1	0	0	1	1	1	0	0
1	0	1	0	0	0	0	0	0	0	1	1
1	0	1	1	1	0	0	0	1	0	0	1
1	1	0	0	0	1	0	1	1	1	0	0
1	1	0	1	0	0	0	1	0	1	1	0
1	1	1	0	0	1	0	0	1	0	0	1
1	1	1	1	0	0	0	0	0	0	1	1

Rising edges are detected when DLY_INDATA is zero, and INDATA is one.   Falling edges when DLY_INDATA is one, and INDATA is zero.   Both edges are detected by simply combining columns.   Neither edge is detected by inverting the "both" column.

5.1 *Example (Edge Detector Initialization)*

```
#define EDGE_RISE0 0x0A0A
#define EDGE_RISE1 0x00CC
#define EDGE_FALL0 0x5050
#define EDGE_FALL1 0x3300
#define EDGE_BOTH0 (EDGE_RISE0|EDGE_FALL0)
#define EDGE_BOTH1 (EDGE_RISE1|EDGE_FALL1)
#define EDGE_NEITHER0 (~EDGE_BOTH0 & 0xFFFF)
#define EDGE_NEITHER1 (~EDGE_BOTH1 & 0xFFFF)

void write_edge (int edgesel,
  unsigned int rising_edge, unsigned int falling_edge, unsigned int neither_edge)
{
  write_select (0x34 + (edgesel&1));

  unsigned int lutvalue=0;
  unsigned int bitmask = 0x80000000;
  for (int i=0; i<16; i=i+1) {
    // Evaluate indata bit1...
    if (neither_edge && bitmask)
      lutvalue |= EDGE_NEITHER1;
    else {
      if (rising_edge & bitmask) lutvalue |= EDGE_RISE1;
      if (falling_edge & bitmask) lutvalue |= EDGE_FALL1;
    }
    bitmask >>= 1;

    // Evaluate indata bit0...
    if (neither_edge && bitmask)
      lutvalue |= EDGE_NEITHER0;
    else {
      if (rising_edge & bitmask) lutvalue |= EDGE_RISE0;
      if (falling_edge & bitmask) lutvalue |= EDGE_FALL0;
    }
    bitmask >>= 1;

    if ((i&1)==0)
      lutvalue <<= 16;
    else {
      write_chain (lutvalue); // write total of 256 bits
      lutvalue = 0;
    }
  }
}
```

Figure 16 - Initialize Edge Detector LUT Values

6   *Timers*

Timer limits are very simple.   You write a 36 bit value.   Lower 32 bits to trigger register 0x38/0x3A. Upper 4 bits to trigger register 0x39/0x3B.   The fpga reference clock is 100Mhz, so each timer tick corresponds to 10ns.
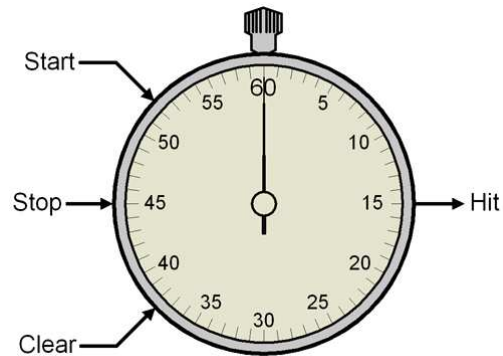
Figure 17 - Timer

The 36-bit timers have a range from 10ns to 687 seconds (over 11 minutes).   Timers are started, stopped, and cleared under control of the trigger sequencer.

6.1   *Example (Timer Initialization)*

```
void write_trigger_limit (int timersel, uint64_t value)
{
  write_select (0x38 + (timersel&1)*2);
  write_chain (value & 0xFFFFFFFF);
  write_select (0x39 + (timersel&1)*2);
  write_chain (value>>32);
}
```

Figure 18 - Initialize Timer Values

7 *Trigger Sums*

The trigger summing logic combines the results of the trigger terms, range detectors, edge detectors and timers.
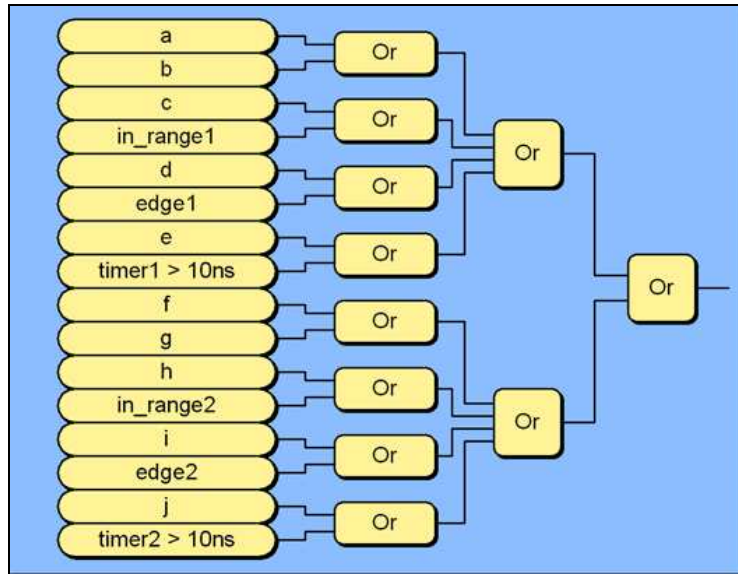
Figure 19 - Combination Trigger Sum

Each of the operations (the "OR" blocks shown above) are implemented using LUT memories. Therefore, they can be configured with any logical operation.   Typical functions are: AND, NAND, OR, NOR, XOR, NXOR, A-only, B-only, ANY, or NOP.

In addition, the LUT memories can be used to invert the result of a trigger term, range, edge or timer detect (anything in the first column above).   For example, allowing "a" or "not a", or "in_range" or "out of range" variations, etc...

There are three trigger sums per sequencer state, and sixteen trigger sequencer states, for a total of 48 trigger sums which need configuring.

7.1 *Trigger Sum Inputs*

As mentioned earlier, the inputs to the sum are trigger terms, range detectors, edge detectors and timers. Results of all inputs feed into the "pair" LUT's (which evaluate a pair of sum inputs each).

- Trigger Terms. 32-bit masked compares. Each term supplies a two bit result. LSB corresponds to indata[15:0]. MSB to indata[31:16].   Both bits must be asserted to indicate a hit.

- Range Detectors. The range detectors similarly supply two signals. LSB is the lower limit hit. MSB is upper limit hit.   Both bits must be asserted to indicate a hit.

- Edge Detectors. Single hit/miss signal.   Fed into both bits 2 & 3 of pair LUT.

- Timer. Single hit/miss signal. Fed into both bits 2 & 3 of pair LUT.

7.2   *Trigger Sum Operations*

Operation codes for the "pair" LUT's are a little confusing since terms & range detectors supply two signals each (both of which must be asserted for a hit).

Sum Inputs				AND	NAND	OR	NOR	XOR	NXOR	A	B
0	0	0	0	0	1	0	1	0	1	0	0
0	0	0	1	0	1	0	1	0	1	0	0
0	0	1	0	0	1	0	1	0	1	0	0
0	0	1	1	0	1	1	0	1	0	1	0
0	1	0	0	0	1	0	1	0	1	0	0
0	1	0	1	0	1	0	1	0	1	0	0
0	1	1	0	0	1	0	1	0	1	0	0
0	1	1	1	0	1	1	0	1	0	1	0
1	0	0	0	0	1	0	1	0	1	0	0
1	0	0	1	0	1	0	1	0	1	0	0
1	0	1	0	0	1	0	1	0	1	0	0
1	0	1	1	0	1	1	0	1	0	1	0
1	1	0	0	0	1	1	0	1	0	0	1
1	1	0	1	0	1	1	0	1	0	0	1
1	1	1	0	0	1	1	0	1	0	0	1
1	1	1	1	1	0	1	0	0	1	1	1

Figure 20 - Pair LUT Operations

The mid LUT is straight forward.   A pure 4-input logic gate:

Sum Inputs				AND	NAND	OR	NOR	XOR	NXOR
0	0	0	0	0	1	0	1	0	1
0	0	0	1	0	1	1	0	1	0
0	0	1	0	0	1	1	0	1	0
0	0	1	1	0	1	1	0	0	1
0	1	0	0	0	1	1	0	1	0
0	1	0	1	0	1	1	0	0	1
0	1	1	0	0	1	1	0	0	1
0	1	1	1	0	1	1	0	0	1
1	0	0	0	0	1	1	0	1	0
1	0	0	1	0	1	1	0	0	1
1	0	1	0	0	1	1	0	0	1
1	0	1	1	0	1	1	0	0	1
1	1	0	0	0	1	1	0	0	1
1	1	0	1	0	1	1	0	0	1
1	1	1	0	0	1	1	0	0	1
1	1	1	1	1	0	1	0	0	1

Figure 21 - Mid LUT Operations

The final LUT only has two inputs. It's configured as a 2-input logic gate:

Sum Inputs				AND	NAND	OR	NOR	XOR	NXOR
0	0	0	0	0	1	0	1	0	1
0	0	0	1	0	1	1	0	1	0
0	0	1	0	0	1	1	0	1	0
0	0	1	1	1	1	1	0	0	1
0	1	0	0	0	0	0	0	0	0
0	1	0	1	0	0	0	0	0	0
0	1	1	0	0	0	0	0	0	0
0	1	1	1	0	0	0	0	0	0
1	0	0	0	0	0	0	0	0	0
1	0	0	1	0	0	0	0	0	0
1	0	1	0	0	0	0	0	0	0
1	0	1	1	0	0	0	0	0	0
1	1	0	0	0	0	0	0	0	0
1	1	0	1	0	0	0	0	0	0
1	1	1	0	0	0	0	0	0	0
1	1	1	1	0	0	0	0	0	0

Figure 22 - Final LUT Operations

All LUT's can also be programmed to NOP (ie: always output zero), or ANY (always output one), or anything else you can think up.

7.3 *Example (Trigger Sum Initialization)*

```
#define OP_NOP 0
#define OP_ANY 1
#define OP_AND 2
#define OP_NAND 3
#define OP_OR 4
#define OP_NOR 5
#define OP_XOR 6
#define OP_NXOR 7
#define OP_A 8
#define OP_B 9

//                  NOP    ANY    AND    NAND   OR     NOR    XOR    NXOR   A      B
int pairvalue[] =  {0x0000,0xFFFF,0x8000,0x7FFF,0xF888,0x0777,0x7888,0x8777,0x8888,0xF000};
int midvalue[] =   {0x0000,0xFFFF,0x8000,0x7FFF,0xFFFE,0x0001,0x0116,0xFEE9,0xEEEE,0xFFF0};
int finalvalue[] = {0x0000,0xFFFF,0x0008,0x0007,0x000E,0x0001,0x0006,0x0009,0x0002,0x0004};

void write_trigger_sum (
  int statenum, int stateterm,
  int op_ab, int op_c_range1, int op_d_edge1, int op_e_timer1,
  int op_fg, int op_h_range2, int op_i_edge2, int op_j_timer2,
  int op_mid1, int op_mid2, int op_final)
{
  write_select (0x40 + (statenum*4) + stateterm);
  write_chain (finalvalue[op_final]);
  write_chain ((midvalue[op_mid2]<<16) | midvalue[op_mid1]);
  write_chain ((pairvalue[op_j_timer2]<<16) | pairvalue[op_i_edge2]);
  write_chain ((pairvalue[op_h_range2]<<16) | pairvalue[op_fg]);
  write_chain ((pairvalue[op_e_timer1]<<16) | pairvalue[op_d_edge1]);
  write_chain ((pairvalue[op_c_range1]<<16) | pairvalue[op_ab]);
}
```

Figure 23 - Initialize Trigger Sum

8   *Trigger Sequence States*

Sequence states allow you to control how the analyzer evaluates and stores sampled data.   Each state gives you control over data capturing, searching for a target (optionally multiple times), or branching to another sequence should a special condition occur.   On a hit, the analyzer advances to the next sequence state. It can also trigger the analyzer controller & start/stop timers. With 16 sequence states, you can create complex protocol following triggers.
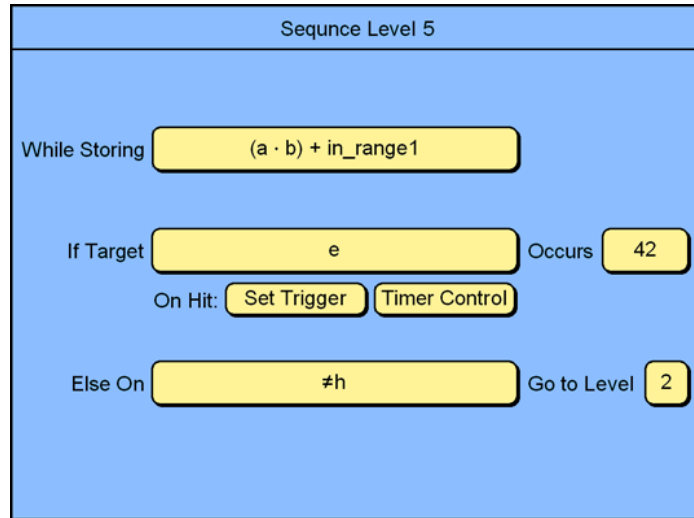
Figure 24 - Example Sequence Level

For each state you specify:

- Capture Term (sum of trigger terms, range checks, edge detect & timer hits)
- Hit Term (sum of trigger terms, range checks, edge detect & timer hits)
- Else Term (sum of trigger terms, range checks, edge detect & timer hits)
- Occurrence Count
- Else State
- Timer Control Flags (start/stop/clear for timers 1 & 2)
- Trigger Flag

The trigger sequencer itself is a relatively simple state machine.   Pseudocode for it is:

```
if (sampled_input valid) then {
    if (capture_term valid) then
        capture data

    if (hit_term valid) then {
        increment hit_count
        if (hit_count == fsm_occurence_count) then {
            update_timers
            reset hit_count
            if (fsm_trigger || last_state) then assert run to controller
            if (!last_state && !fsm_last_state) advance to next state
        }
    }

    otherwise if (else_term valid) then {
        reset hit_count
        go to state "fsm_else_level"
    }
}
```

Figure 25 - Trigger Sequencer State Machine

8.1   *Example (Trigger Sequence Initialization)*

```
#define TRIGSTATE_STATENUM_MASK 0xF
#define TRIGSTATE_OBTAIN_MASK   0x000FFFFF
#define TRIGSTATE_ELSE_BITOFS   20
#define TRIGSTATE_STOP_TIMER0   0x01000000
#define TRIGSTATE_STOP_TIMER1   0x02000000
#define TRIGSTATE_CLEAR_TIMER0  0x04000000
#define TRIGSTATE_CLEAR_TIMER1  0x08000000
#define TRIGSTATE_START_TIMER0  0x10000000
#define TRIGSTATE_START_TIMER1  0x20000000
#define TRIGSTATE_TRIGGER_FLAG  0x40000000
#define TRIGSTATE_LASTSTATE     0x80000000

void write_trigger_state (
  int statenum,      // 0 to 15
  int last_state,
  int set_trigger,
  int start_timer,   // bit0=timer1, bit1=timer2
  int stop_timer,    // bit0=timer1, bit1=timer2
  int clear_timer,   // bit0=timer1, bit1=timer2
  int else_state,    // 0 to 15
  int obtain_count)
{
  write_select (statenum & TRIGSTATE_STATENUM_MASK);

  unsigned int value =
    ((else_state & TRIGSTATE_STATENUM_MASK)<<TRIGSTATE_ELSE_BITOFS) |
    (obtain_count & TRIGSTATE_OBTAIN_MASK);
  if (last_state) value |= TRIGGER_LASTSTATE;
  if (set_trigger) value |= TRIGSTATE_TRIGGER_FLAG;
  if (start_timer & 1) value |= TRIGSTATE_START_TIMER0;
  if (start_timer & 2) value |= TRIGSTATE_START_TIMER1;
  if (stop_timer & 1) value |= TRIGSTATE_STOP_TIMER0;
  if (stop_timer & 2) value |= TRIGSTATE_STOP_TIMER1;
  if (clear_timer & 1) value |= TRIGSTATE_CLEAR_TIMER0;
  if (clear_timer & 2) value |= TRIGSTATE_CLEAR_TIMER1;
  write_chain (value);
}
```

Figure 26 - Initialize Trigger Sequence

Appendix A - Advanced Trigger Config Registers

The client initializes advanced trigger LUT chains using the 0x9E *configuration-select* command, followed by one or more 0x9F *write-data* commands. Each data write adds 32-bits to the selected LUT chain.

Exceptions are the FSM state configs & timer limits, which store DWORD values directly. However, the client sees no difference (just more writes to trigger config regs).

Client Long Commands (command byte + four data bytes, LSB data byte first):
```
  0x9E = ADVTRIG-Config-Select (LSB = Chain#.   Other bytes reserved.)
  0x9F = ADVTRIG-Write-Data
```

FSM State Data
```
  0x00      = FSM State 0      (32-bit word's)
  0x01      = FSM State 1
  0x02      = FSM State 2
  0x03      = FSM State 3
  0x04      = FSM State 4
  0x05      = FSM State 5
  0x06      = FSM State 6
  0x07      = FSM State 7
  0x08      = FSM State 8
  0x09      = FSM State 9
  0x0A      = FSM State 10
  0x0B      = FSM State 11
  0x0C      = FSM State 12
  0x0D      = FSM State 13
  0x0E      = FSM State 14
  0x0F      = FSM State 15
```

LUTCHAIN's:
```
  0x20      = trigger term a  (128 bit chains)
  0x21      = trigger term b
  0x22      = trigger term c
  0x23      = trigger term d
  0x24      = trigger term e
  0x25      = trigger term f
  0x26      = trigger term g
  0x27      = trigger term h
  0x28      = trigger term i
  0x29      = trigger term j
  0x30      = range 1 lower   (512 bit chains)
  0x31      = range 1 upper   (512 bit chains)
  0x32      = range 2 lower   (512 bit chains)
  0x33      = range 2 upper   (512 bit chains)
  0x34      = edge 1          (64 bit chains)
  0x35      = edge 2          (64 bit chains)
  0x38-0x39 = timer 1 limit   (36-bit word's)
  0x3A-0x3B = timer 2 limit   (36-bit word's)

  0x40      = state 0 hit-term
  0x41      = state 0 else-term
  0x42      = state 0 capture-term
  0x44-0x46 = state 1 terms
  0x48-0x4A = state 2 terms
  0x4C-0x4E = state 3 terms
  0x50-0x52 = state 4 terms
  0x54-0x56 = state 5 terms
  0x58-0x5A = state 6 terms
  0x5C-0x5E = state 7 terms
  0x60-0x62 = state 8 terms
  0x64-0x66 = state 9 terms
  0x68-0x6A = state 10 terms
  0x6C-0x6E = state 11 terms
  0x70-0x72 = state 12 terms
  0x74-0x76 = state 13 terms
  0x78-0x7A = state 14 terms
  0x7C-0x7E = state 15 terms
```

Appendix B - Example Code

```
#define write_select(value) write_long_command (0x9E, value)
#define write_chain(value) write_long_command (0x9F, value)

#define RANGE_XOR0 0xAAAA
#define RANGE_XOR1 0x5555
#define RANGE_NOP 0xFFFF

#define EDGE_RISE0 0x0A0A
#define EDGE_RISE1 0x00CC
#define EDGE_FALL0 0x5050
#define EDGE_FALL1 0x3300
#define EDGE_BOTH0 (EDGE_RISE0|EDGE_FALL0)
#define EDGE_BOTH1 (EDGE_RISE1|EDGE_FALL1)
#define EDGE_NEITHER0 (~EDGE_BOTH0 & 0xFFFF)
#define EDGE_NEITHER1 (~EDGE_BOTH1 & 0xFFFF)

#define TRIGSTATE_STATENUM_MASK 0xF
#define TRIGSTATE_OBTAIN_MASK   0x000FFFFF
#define TRIGSTATE_ELSE_BITOFS   20
#define TRIGSTATE_STOP_TIMER0   0x01000000
#define TRIGSTATE_STOP_TIMER1   0x02000000
#define TRIGSTATE_CLEAR_TIMER0 0x04000000
#define TRIGSTATE_CLEAR_TIMER1 0x08000000
#define TRIGSTATE_START_TIMER0 0x10000000
#define TRIGSTATE_START_TIMER1 0x20000000
#define TRIGSTATE_TRIGGER_FLAG 0x40000000
#define TRIGSTATE_LASTSTATE     0x80000000

#define OP_NOP 0
#define OP_ANY 1
#define OP_AND 2
#define OP_NAND 3
#define OP_OR 4
#define OP_NOR 5
#define OP_XOR 6
#define OP_NXOR 7
#define OP_A 8
#define OP_B 9


//
// Setup LUT's for Trigger Term
//
// Inputs are 32-bit target & mask for comparing against captured analyzer data.
// If a mask bit is set, the corresponding target bit participates in the trigger.
//
void write_term (int term_number, unsigned int target, unsigned int mask)
{
  unsigned int bitmask=1;
  unsigned int lutvalue0=0;
  unsigned int lutvalue1=0;
  unsigned int lutvalue2=0;
  unsigned int lutvalue3=0;

  for (int i=0; i<16; i=i+1) {
    if (((i ^ (target & 0xF)) & (mask & 0xF))==0) lutvalue0 |= bitmask;
    if (((i ^ ((target>>4) & 0xF)) & ((mask>>4) & 0xF))==0) lutvalue0 |= (bitmask<<16);
    if (((i ^ ((target>>8) & 0xF)) & ((mask>>8) & 0xF))==0) lutvalue1 |= bitmask;
    if (((i ^ ((target>>12) & 0xF)) & ((mask>>12) & 0xF))==0) lutvalue1 |= (bitmask<<16);
    if (((i ^ ((target>>16) & 0xF)) & ((mask>>16) & 0xF))==0) lutvalue2 |= bitmask;
    if (((i ^ ((target>>20) & 0xF)) & ((mask>>20) & 0xF))==0) lutvalue2 |= (bitmask<<16);
    if (((i ^ ((target>>24) & 0xF)) & ((mask>>24) & 0xF))==0) lutvalue3 |= bitmask;
    if (((i ^ ((target>>28) & 0xF)) & ((mask>>28) & 0xF))==0) lutvalue3 |= (bitmask<<16);
    bitmask <<= 1;
  }

  // Write data into LUT serial chain.  MSB must goes in first.  Total of 128 bits.
  write_select (0x20 + (term_number%10));
  write_chain (lutvalue3);
  write_chain (lutvalue2);
  write_chain (lutvalue1);
  write_chain (lutvalue0);
}
```

```
//
// Setup LUT's for Range Detectors
//
// Inputs:
//   rangesel:  0=range1-lower, 1=range1-upper, 2=range2-lower, 3=range2-upper
//   target:    Limit value.
//   mask:      Indicate which bits should participate in range compare.
//
void write_range (int rangesel, unsigned int target, unsigned int mask)
{
  unsigned int value;
  unsigned int lutvalue=0;
  int i;

  write_select(0x30+(rangesel&3));

  // Count # of bits in mask...
  int bitcount=0;
  unsigned int bitmask=1;
  for (i=0; i<32; i++, bitmask<<=1)
    if (mask & bitmask)
      bitcount++;

  // Prepare target value...
  if (rangesel & 1)
    value = ~target; // upper range target
  else value = ~(target-1); // lower range value

  // Push MSB of target into bit 31...
  value <<= (32-bitcount);

  // Generate & program LUT values.  Total of 512 bits.
  for (i=0; i<16; i=i+1) {
    if (((mask>>31)&1)==0)
      lutvalue = RANGE_NOP;
    else {
      lutvalue = ((value>>31)&1) ? RANGE_XOR1 : RANGE_XOR0;
      value <<= 1;
    }
    mask <<= 1;
    lutvalue <<= 16;

    if (((mask>>31)&1)==0)
      lutvalue |= RANGE_NOP;
    else {
      lutvalue |= ((value>>31)&1) ? RANGE_XOR1 : RANGE_XOR0;
      value <<= 1;
    }
    mask <<= 1;

    write_chain(lutvalue);
  }
}


//
// Setup LUT's for Edge Detectors.
// Each LUT handles two bits of input.  Bits 0 & 1 are input.  Bits 2 & 3 delayed input.
//
void write_edge (int edgesel,
  unsigned int rising_edge, unsigned int falling_edge, unsigned int neither_edge)
{
  write_select (0x34 + (edgesel&1));

  unsigned int lutvalue=0;
  unsigned int bitmask = 0x80000000;
  for (int i=0; i<16; i=i+1) {
    // Evaluate indata bit1...
    if (neither_edge && bitmask)
      lutvalue |= EDGE_NEITHER1;
    else {
      if (rising_edge & bitmask) lutvalue |= EDGE_RISE1;
      if (falling_edge & bitmask) lutvalue |= EDGE_FALL1;
    }
    bitmask >>= 1;

    // Evaluate indata bit0...
    if (neither_edge && bitmask)
      lutvalue |= EDGE_NEITHER0;
    else {
      if (rising_edge & bitmask) lutvalue |= EDGE_RISE0;
```

```
      if (falling_edge & bitmask) lutvalue |= EDGE_FALL0;
    }
    bitmask >>= 1;

    if ((i&1)==0)
      lutvalue <<= 16;
    else {
      write_chain (lutvalue); // write total of 256 bits
      lutvalue = 0;
    }
  }
 }
}


//
// Setup trigger timers...
//
void write_trigger_limit (int timersel, uint64_t value)
{
  write_select (0x38 + (timersel&1)*2);
  write_chain (value & 0xFFFFFFFF);
  write_select (0x39 + (timersel&1)*2);
  write_chain (value>>32);
}


//
// Setup trigger state term combinational sum.    Operations for
// all 11 fields are merged & written.  Uses table based lookups.
//
//                    NOP    ANY    AND    NAND   OR     NOR    XOR    NXOR   A      B
int pairvalue[] =  {0x0000,0xFFFF,0x8000,0x7FFF,0xF888,0x0777,0x7888,0x8777,0x8888,0xF000};
int midvalue[] =   {0x0000,0xFFFF,0x8000,0x7FFF,0xFFFE,0x0001,0x0116,0xFEE9,0xEEEE,0xFFF0};
int finalvalue[] = {0x0000,0xFFFF,0x0008,0x0007,0x000E,0x0001,0x0006,0x0009,0x0002,0x0004};

void write_trigger_sum (
  int statenum, int stateterm,
  int op_ab, int op_c_range1, int op_d_edge1, int op_e_timer1,
  int op_fg, int op_h_range2, int op_i_edge2, int op_j_timer2,
  int op_mid1, int op_mid2, int op_final)
{
  write_select (0x40 + (statenum*4) + stateterm);
  write_chain (finalvalue[op_final]);
  write_chain ((midvalue[op_mid2]<<16) | midvalue[op_mid1]);
  write_chain ((pairvalue[op_j_timer2]<<16) | pairvalue[op_i_edge2]);
  write_chain ((pairvalue[op_h_range2]<<16) | pairvalue[op_fg]);
  write_chain ((pairvalue[op_e_timer1]<<16) | pairvalue[op_d_edge1]);
  write_chain ((pairvalue[op_c_range1]<<16) | pairvalue[op_ab]);
}


//
// Setup trigger state.    Specify trigger mode, timer controls, "else" state number,
// and necessary "obtain" count (number of hits needed before full hit observed).
//
void write_trigger_state (
  int statenum,      // 0 to 15
  int last_state,
  int set_trigger,
  int start_timer,   // bit0=timer1, bit1=timer2
  int stop_timer,    // bit0=timer1, bit1=timer2
  int clear_timer,   // bit0=timer1, bit1=timer2
  int else_state,    // 0 to 15
  int obtain_count)
{
  write_select (statenum & TRIGSTATE_STATENUM_MASK);

  unsigned int value =
    ((else_state & TRIGSTATE_STATENUM_MASK)<<TRIGSTATE_ELSE_BITOFS) |
    (obtain_count & TRIGSTATE_OBTAIN_MASK);
  if (last_state) value |= TRIGSTATE_LASTSTATE;
  if (set_trigger) value |= TRIGSTATE_TRIGGER_FLAG;
  if (start_timer & 1) value |= TRIGSTATE_START_TIMER0;
  if (start_timer & 2) value |= TRIGSTATE_START_TIMER1;
  if (stop_timer & 1) value |= TRIGSTATE_STOP_TIMER0;
  if (stop_timer & 2) value |= TRIGSTATE_STOP_TIMER1;
  if (clear_timer & 1) value |= TRIGSTATE_CLEAR_TIMER0;
  if (clear_timer & 2) value |= TRIGSTATE_CLEAR_TIMER1;
  write_chain (value);
}
```

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